

We would like to compute the powers (resp. the exponential or any matrix function) of an  $n \times n$  matrix  $A$ .

➤ **Case 1**

The matrix  $A$  is **diagonalizable**:

$$A_m(\lambda_i) = G_m(\lambda_i) \text{ for } 1 \leq i \leq k.$$

We need to apply **Cayley-Hamilton Theorem**:

➔  $v$  is called an eigenvector for  $\lambda$  if  $(A - \lambda I)v = 0$  with  $v \neq 0$ .

In both cases we must first compute the characteristic polynomial and determine the eigenvalues and eigenvectors.

• We also use **Minimal polynomial**.



➤ **Case 2**

The matrix  $A$  is **not diagonalizable**.

We need to apply **Cayley-Hamilton Theorem**:  $p_A(A) = 0$

In this case we must compute the **generalized eigenvectors** of  $A$ .

➔  $v$  is called a generalized eigenvector for  $\lambda$  if  $(A - \lambda I)^s v = 0$  for some  $s \geq 2$

**Theorem (Necessary and sufficient condition for diagonalization)**

Let  $A \in \mathcal{M}_n(\mathbb{K})$  be a square matrix.  $A$  is diagonalizable, if and only if, there exists a basis  $\mathcal{B}$  of  $\mathbb{K}^n$  formed by  $n$  eigenvectors of  $A$ .

• If  $A = P \cdot D \cdot P^{-1}$ , then  $f(A) = P \cdot f(D) \cdot P^{-1}$ .

Here,  $P$  is invertible and  $D$  is diagonal.

**Example**

Assume that  $A = PDP^{-1}$ , where  $P$  is invertible and  $D$  is diagonal. We have

- For  $f(x) = x^k \Rightarrow f(A) = A^k = P \cdot D^k \cdot P^{-1}$  for  $k \geq 0$
- For  $f(x) = \sqrt{x} \Rightarrow f(A) = \sqrt{A} = P \cdot \sqrt{D} \cdot P^{-1}$
- For  $f(x) = \cos x \Rightarrow f(A) = \cos A = P \cdot \cos D \cdot P^{-1}$
- For  $f(x) = e^x \Rightarrow f(A) = e^A = P \cdot e^D \cdot P^{-1}$
- For  $f(x) = \ln x \Rightarrow f(A) = \ln A = P \cdot \ln D \cdot P^{-1}$
- ...
- and so on.

➤ We apply the powers of the matrix  $A$  in solving **systems of recurrence sequences** and its exponential in solving **linear systems of differential equations** of the first order.

•  $X_{n+1} = AX_n \Rightarrow X_n = A^n X_0$ .

•  $X_{n+1} = AX_n + C \Rightarrow X_n = A^n X_0 + (A^{n-1} + A^{n-2} + \dots + A + I)C$ .

•  $X' = AX \Rightarrow X(t) = c_1 e^{\lambda_1 t} v_1 + c_2 e^{\lambda_2 t} v_2 + \dots + c_n e^{\lambda_n t} v_n$

**Conclusion.** We can compute the powers and exponential of any  $n$  by  $n$  matrix.

• If  $A = P \cdot J \cdot P^{-1}$ , then  $e^{At} = P \cdot e^{Jt} \cdot P^{-1}$ .

←  $J^k$  and  $e^J$  ←

We have the formulas of the powers and exp of  $J$

➤ We will apply one of the following three famous Theorems:

1. **Schur decomposition Theorem**: States that every  $n$  by  $n$  matrix is similar to an upper triangular matrix  $T$ .
2. **Dunford Decomposition Theorem**: States that every  $n$  by  $n$  matrix  $A$  can be written as  $A = D + N$ , where  $D$  is **diagonalizable**,  $N$  is **nilpotent** with  $D \cdot N = N \cdot D$ . Here, for computing the powers of  $A$ , we must use the **Binomial formula** for matrices:

$$A^n = (D + N)^n = C_n^0 D^n + C_n^1 D^{n-1} N + \dots + C_n^{k-1} D^{n-(k-1)} N^{k-1},$$

where  $N^k = 0$ .

3. **Jordan Decomposition Theorem**: States that every  $n$  by  $n$  matrix  $A$  is similar to a Jordan matrix by **blocks**, i.e.

where  $J$  is given by:

$$A = P \cdot J \cdot P^{-1}$$

$$J = \begin{pmatrix} J_{\lambda_1} & & & \\ & J_{\lambda_2} & & \\ & & \ddots & \\ & & & J_{\lambda_k} \end{pmatrix}$$

with

$$J_{\lambda_i} = \begin{pmatrix} \lambda_i & 1 & & \\ & \lambda_i & 1 & \\ & & \ddots & \ddots \\ & & & \lambda_i \end{pmatrix} \text{ or } \begin{pmatrix} \lambda_i & & & \\ 1 & \lambda_i & & \\ & 1 & \ddots & \\ & & & \ddots & \lambda_i \end{pmatrix}$$